

**SIGNETICS TWIN
2650
RELOCATABLE
ASSEMBLER
MANUAL**



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SIGNETICS TWIN 2650 RELOCATABLE ASSEMBLER MANUAL

signetics

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FOREWORD

The TWIN Relocatable Assembler (RASM) provides features, designed for the development of larger software systems, which are not available with the TWIN Absolute Assembler (ASM). For RASM version 3.0, the list of features include:

1. A relocatable output is produced, so that locations in which a program is to be executed can be determined after assembling.
2. It is possible to introduce items which are defined in other programs so links between separately assembled programs can be established mechanically.
3. The mnemonic branch-instruction codes are extended by the inclusion of the condition in the mnemonic, increasing the readability of programs. For example, BER is equivalent to BCTR,0.
4. ASCII data constants and other types of data constants may be combined into a single statement.
5. Data definitions may be preceded by a multiplication factor.
6. The number of bytes in a character string (ASCII data constant) can be specified explicitly.
7. The new assembler accepts most correct ASM programs. Major differences are in the ORG and END definitions.

In addition to RASM, a link editor, called LINK, has been constructed to provide the capability of linking separately assembled programs. RASM produces an object module and LINK links a number of these object modules and edits the modules into one or more load modules. The modules can be loaded by means of the SDOS commands LOAD/GO and XEQ.

NOTE

The object file output of RASM is not compatible with the hexadecimal output of RASM. The WHEX command can be used to reformat LINK output into hexadecimal format.

RASM assembles source code for the 2650B microprocessor and because the 2650A instruction set is a proper subset of 2650B assembly language, 2650A code may also be assembled with RASM.

CAUTION

The 2650B instructions STPL and LDPL are not members of the 2650A instruction set. When these instructions appear in the program, the user is required to know on which CPU the program is to execute.



I. INTRODUCTION

The 2650 assembly language is a symbolic language which facilitates the writing of programs for the Signetics 2650 microprocessor. The 2650 Relocatable Assembler, RASM, is a program which accepts this symbolic source code as input and produces a listing and/or a relocatable program object module as output.

The assembler is a two-pass program that builds a symbol table, issues helpful error messages and produces an easily readable program listing. In conjunction with the link editor, LINK, a computer readable object (load) module is output. It features conditional assembly, symbolic and relative addressing, forward references, free format source code, self-defining constants, complex expression evaluation, relocatability, and a versatile set of Pseudo Operations. Additionally, the assembler is capable of generating data in several number-based systems, and ASCII character code. These features aid the programmer/engineer in producing well documented, working programs in minimum time.

ASSEMBLY LANGUAGE

An assembly language program is a program written in symbolic machine language. It is comprised of statements. A statement is either a symbolic machine instruction, a pseudo-operation statement, or a comment.

A symbolic machine instruction is a written specification for a particular machine operation expressed by a symbolic operation code and, if required, a symbolic address or operand. For example:

```
LOC2    STRR,RO    SAV
```

Where:

LOC2	is a symbol representing the memory address of the instruction.
STRR	is a symbolic operation code representing the bit pattern of the "store relative" instruction.
RO	is a symbol that has been defined as register 0 by the "EQU" pseudo-op.
SAV	is a symbol representing the memory location into which the contents of register 0 are to be stored.

A pseudo-operation statement is a statement which is not translated into a machine instruction, but rather is interpreted as a directive to the assembler program. For example:

```
SCHD    ACON    REDY
```

Where:

SCHD	is a symbol. The assembler is to assign the memory address of the first byte of the two allocated to this symbol.
ACON	is a pseudo-op which directs the assembler program to allocate two bytes of memory.
REDY	is a symbol representing an address. The assembler is directed to place the equivalent memory address into the allocated bytes.

STATEMENTS

Statements are always written in a particular format. The format is depicted below:

LABEL FIELD	OPERATION FIELD	OPERAND FIELD	COMMENT FIELD
-------------	-----------------	---------------	---------------

The statement is always assumed to be written as an 80-column image.

The Label Field is provided to assign symbolic names to bytes of memory. If present, the Label Field must begin in column one.

The Operation Field is provided to specify a symbolic operation code or a pseudo-operation code. If present, the Operation Field must either begin past column one or be separated from the Label Field by one or more blanks.

The Operand Field is provided to specify arguments for the operation in the Operation Field. The Operand Field, if present, is separated from the Operation Field by one or more blanks.

The Comment Field is provided to enable the assembly language programmer optionally to place a message stating the purpose or intent of a statement or a group of statements. The Comment Field must be separated from the preceding field by one or more blanks.

Comment Statement

A Comment Statement is a statement that is not processed by the assembler program. It is merely reproduced on the assembly listing. A Comment Statement is indicated by encoding an asterisk in column one. For example:

*THIS IS A COMMENT STATEMENT

LOCATION COUNTER

During the assembly process, the assembler maintains a cell that always contains the address of the next memory location to be assembled. This cell is called the Location Counter. It is used by the assembler to assign addresses to assembled bytes, and its value is available to be used by the programmer.

The character "\$" is a symbol that the assembler recognizes as the symbolic name of the Location Counter. It may be used like any other symbol, but it may not appear in the label field.

When using the "\$," the programmer may think of it as expressing the idea "\$" = "address of myself." For example:

```
10816    BCTR,3    $
```

The first byte of this branch instruction is in location 108₁₆. The instruction directs the microprocessor to "branch to myself." The Location Counter in this example contains the value 108₁₆.

SYMBOLIC ADDRESSING

As mentioned above, the user can attach a label to an instruction, as shown in the following example:

```
SAVR    STRR,RO    SAV
```

The assembler, upon seeing a valid symbol in the label field, assigns the equivalent address (the value of the Location Counter) to the label. In the given example, if the STRR instruction is to be stored in the address 0127, then the symbol SAVR would be made equivalent to the value 0127 for the duration of the assembly.

The symbol could then be used anywhere in the source program to refer to the address value or, more typically, it could be used to refer to the instruction location. The important concept is that the address of the instruction need not be known; only the symbol need be used to refer to the instruction location. Thus, when branching to the STRR instruction, one could write:

```
BCTA,3    SAVR
```

When this three-byte branch instruction is translated by the assembler, the address of the STRR instruction is placed in the address field of the branch instruction. Similarly, in the STRR instruction above, the symbol SAV in the Operand Field is replaced by the address of the memory location whose label is SAV.

It is also possible to use symbolic addresses which are near other locations to refer to those locations without defining new labels. For example:

	BCTR,3	BEG
	BCTR,0	BEG+4
	ANDZ	3
	BSTR,3	S+48
BEG	LODA,2	PAL
	HALT	
	SUBI,2	3

In the above example, the instruction "BCTR,3 BEG" refers to the "LODA,2 PAL" instruction. The instruction "BCTR,0 BEG+4" refers to the "SUBI,2 3" instruction.

BEG+4 means the address BEG plus four bytes. This type of expression is called relative symbolic addressing and given a symbolic address; it can be used as a landmark to express several bytes before or after the symbolic address. Examples are:

	BCTR,3	PAL+23
	BSTA,0	STT-18

The arguments are evaluated like expressions and cannot exceed in value the maximum number that can be contained in an integer constant in the range of -32,768 to +32,767.

II. LANGUAGE ELEMENTS

Input to the assembler consists of a sequence of characters combined to form assembly language elements. These language elements include symbols, instruction mnemonics, constants and expressions which make up the individual program statements that comprise a source program.

CHARACTERS

Alphabetic: A through Z

Numeric: 0 through 9

Special characters:

	blank
(left parenthesis
)	right parenthesis
+	add or positive value
-	subtract or negative value
*	asterisk
'	single quote
,	comma
/	slash
\$	dollar sign
<	less than sign
>	greater than sign
#	pound sign
@	at sign
?	question mark
!	exclamation point
"	double quote
%	percent sign
=	equal
;	semicolon
:	colon
.	period

SYMBOLS

Symbols are formed from combinations of characters. Symbols provide a convenient means of identifying program elements so that they can be referenced by other elements.

1. Symbols may consist of 1 to 6 characters. Any of the alphabetic, numeric, or the special characters #,@,!,", and % are valid.
2. The character \$ is a special symbol which may be used in the argument field of a statement to represent the current value of the Location Counter.

3. The character * is a special symbol which is used as an indirect address indicator.
4. The characters + and - are also used as auto-increment/auto-decrement indicators.

The following are examples of valid symbols:

DOP1	RAV3	SEVEN%
AA	TEMZ	AT#12

The following are examples of invalid symbols:

ROUND OFF	more than six characters
1LAR	begins with a numeric
PA N	imbedded blank
TEN\$	unallowed special character

CONSTANTS

A constant is a self-defining language element. Unlike a symbol, the value of a constant is its own "face" value and is invariant. Internal numbers are represented in 2's complement notation. There are two forms in which constants may be written: the Self-Defining Constant and the General Constant.

Self-Defining Constant

The self-defining constant is a form of constant which is written directly in an instruction and defines a decimal value. For example:

```
LODA,R3    BUFF+65
```

In this example, 65 is a self-defining constant. The value of the integer constant expressed by a self-defining constant must be in the range -32,768 to +32,767.

General Constant

The general constant is also written directly in an instruction, but the interpretation of its value is dictated by a code character and delimited by quotation marks.

```
LODA,R3    BUFF+H'3E'
```

In this example, the code letter H specifies that 3E is a hexadecimal constant equivalent to decimal value 62.

The size of a number generated by a general constant form (B, O, D, H) must be in the range -32,768 to +32,767. However, the most important concept to

understand when using constant forms is that the final value of a resolved expression must fit the constraints of the actual field destined to contain the value. For example:

```
LODA,R2    PAL+H'3EE2"-H'3EE0'
```

In this case, the argument, when resolved, must fit into the 13 bits in the actual machine instruction. Even though each of the two hexadecimal constants is larger than can fit into 13 bits, the final value of the expression is containable in 13 bits and therefore the constants are permitted. Similarly, the statement `DATA H'3FE'` is not allowed, as the `DATA` statement defines one byte quantities and `H'3FE'` specifies more than 8 bits. Summarily, the size of the evaluated expressions must be less than or equal to their corresponding data fields.

There are five types of General Constants usable with 2650 Assembler:

Code	Type
B	Binary Constant
O	Octal Constant
D	Decimal Constant
H	Hexadecimal Constant
A	ASCII Character Constant

B: Binary Constant

A binary constant consists of an optionally signed binary number of up to 16 bits enclosed in single quotes and preceded by the letter B, e.g., `B'1011011'`. Binary information is stored right justified.

O: Octal Constant

An octal constant consists of an optionally signed octal number enclosed by single quotation marks and preceded by the letter O, e.g., `O'352'`. The value will be right justified.

D: Decimal Constant

A decimal constant consists of an optionally signed decimal number enclosed by single quotation marks and preceded by the letter D, e.g., `D'249'`. The value will be right justified.

H: Hexadecimal Constant

A hexadecimal constant consists of an optionally signed hexadecimal number enclosed in single quotation marks and preceded by the letter H, e.g., `H'3F'`. The value will be right justified.

A: ASCII Character Constant

An ASCII character constant consists of a string of ASCII characters enclosed by single quotation marks and preceded by the letter A. For example, `A'HELLO THERE'`. Each character will be encoded in 7-bit ASCII and stored in

successive bytes. The high-order bit is always set to zero in each allocated byte.

See Appendix D for permissible characters and their equivalent ASCII codes. To specify a single quotation mark as a character constant, it must appear twice in the character string, e.g., A'TYPE"HELP"NOW' will appear in storage as TYPE'HELP'NOW.

MULTIPLE CONSTANT SPECIFICATIONS

General constant forms, when used in DATA and ACON statements, allow multiple specifications within the constant expression. For example, D'52,21,208,27'. A comma separates each byte specification (except in the ASCII form) and successive specifications determine successive bytes of storage. The forms within a single constant expression may be mixed. Each byte may be optionally signed. For example:

```
H'03,-F2,+11,-8,33,0',0'271,133',255,H'F0,FF'0'271,133',A'X,H'0,A,B,C'
```

EXPRESSIONS

An expression is an assembly language element that represents a value. It consists of a single term or a combination of terms separated by arithmetic, logical or relational operators. A term may be a valid symbolic reference, a self-defining constant or a general constant.

The valid operators which may be used in an expression are:

+	for addition (or unary positive sign)
-	for subtraction (or unary negative sign)
*	for multiplication
/	for division
.MOD.	for remainder after division
.NOT.	for one's complement (unary operator)
.AND.	for logical AND
.OR.	for logic OR
.XOR.	for logical EXCLUSIVE OR
.SHR.	for logical shift right
.SHL.	for logical shift left
>	for retrieving the lower byte of a two-byte argument
<	for retrieving the upper byte of a two-byte argument
.EQ.	for equal (=) relation
.NE.	for not equal (\neq) relation
.GT.	for greater than (>) relation
.LT.	for less than (<) relation
.GE.	for greater than or equal to (\geq) relation
.LE.	for less than or equal to (\leq) relation

The relational operators are binary operators which yield a true (1) or false (0) value.

These operators have an implied priority which determines the order in which the operations are performed in multiple operator expressions. The operations specified by the operators of the highest priority are performed first, then those of the second priority, and so forth.

The priority of operators is as follows, with operators on the same line having equal priority:

```
* / .MOD. .SHR. .SHL.  
+ -  
.EQ. .NE. .GT. .LT. .GE. .LE.  
.NOT.  
.AND.  
.OR. .XOR.  
< >
```

Parentheses can be used to override the implied order of operations or to make the expression clearer. For example:

```
A + B * C          is equivalent to A + (B * C)  
.NOT.A-B.MOD.C     is equivalent to .NOT.(A-(B.MOD.C))  
A.AND.B-C          is equivalent to A.AND.(B-C)
```

The expression A.SHR.B shifts the argument A B places towards the right and fills the most-significant bits of A with zeroes. Similarly, A.SHL.B shifts A B places to the left and fills the least-significant bits with zeroes. For example:

```
H'43'.SHR.1        is equivalent to H'21'  
'H'21'.SHL.1      is equivalent to H'42'
```

Examples of valid expressions are:

```
LOOP              PAL-$  
LOOP+5           $-PAL+3  
SAM+3-LOOP       BIT-3+H'3B'
```

NOTE

The special symbol '\$' represents the current value of the Location Counter.

If an expression resolves to a two-byte value where only a single-byte value is required, the assembler will use the least-significant byte as the operand.

EVALUATION OF EXPRESSIONS

Due to the relocation facility and the ability to refer to external symbols, several types of operands in expressions can be distinguished:

1. Constant - A constant is an internally defined item with an invariant value. Constants can be defined by the assembler instruction EQU; for example:

CR EQU H'OD'

2. Address - An address is also an internally defined item, but its final value is assigned by the link editor when the program object code is located. Labels of machine instructions are examples of addresses.
3. External Reference - An external reference is a reference to an item which is defined in another program, where it is either an address or a constant. External references are introduced with the instruction EXTRN.

All operations are allowed on constants, but only a limited number of operations are useful on addresses and external references. When an operation, not giving a usable result, is performed on an address or an external reference, the result is a constant. The operations, types of operands and results are:

Where: A = address; E = external reference; C = constant.

<u>Operation</u>	<u>Operands</u>	<u>Result</u>
Addition:	A + C	A
	E + C	E
		C in all other cases
Subtraction:	A - C	A
	E - C	E
		C in all other cases
Any Other Operation:		C in all cases

Expressions are evaluated with 16-bit arithmetic, that is, division may have an 8-bit divisor and 16-bit dividend; multiplication may have a 16-bit multiplier and multiplicand. When an expression evaluates to larger than a 16-bit signed integer, an overflow error is reported and processing on the line aborts. As for constants, the final value of a resolved expression must fit the constraints of the actual field destined to contain the value.

SPECIAL OPERATORS

There are two special operators that are recognized by the assembler. They are:

< less than sign
> greater than sign

The assembler interprets these operators in a special way:

> perform a modulo 256 divide (use low-order byte)
< perform an integer divide by 256 (use high-order byte)

These special operators are intended to be used to access a two-byte address in one byte parts using a minimum of storage. For example, if it is desired to get the high-order bits of an address (ADDB) into register 2 and the low-order bits into register 1 it could be done as follows:

```
                LODR,R2      APAL  
                LODR,R1      APAL+1  
                . . .  
                . . .  
                . . .  
                ACON          ADDB  
APAL
```

or by utilizing the special operators, it could be done as follows:

```
                LODI,R2      < ADDB  
                LODI,R1      > ADDB
```

The first method uses 6 bytes to accomplish what the second method can do in 4 bytes.

The special operators are most often used to facilitate the passing of an address in registers.

III. SYNTAX

Assembly language elements may be combined to express symbolically both 2650 instructions and assembler directives. There are specific rules for writing these instructions. This set of rules is known as the Syntax of the symbolic assembly language.

FIELDS

A statement prepared for processing by the assembler is divided into four fields: the Label Field, the Operation Field, the Operand Field and the Comment Field. Each field is separated by at least one blank character. Columns 1 through 80 of the input line are scanned by the assembler.

Label Field

The label field optionally contains a symbolic name which the assembler assigns to the instruction specified in the remaining part of the line. If a name is specified, it must begin in column 1. The assembler assumes that there is no label if column 1 is blank. The label field, if present, must contain only a valid symbol.

Operation Field

The operation field contains a mnemonic code which represents a 2650 processor operation or an assembly directive. The operation field must be present in every non-comment line. See Appendix B for a list of the valid mnemonic codes. Additionally, depending on the instruction type, the operation field may also specify a general purpose register or a condition code.

Operand Field

The operand (or argument) field contains one or more symbols, constants or expressions separated by commas. The argument field specifies storage locations, constants, register specifications and any other information necessary to specify completely a machine operation or an assembler directive. Embedded blanks are not permitted as they are considered field terminators.

Comment Field

The comment field contains any valid characters in any combination. The comment field is not processed by the assembler, but is merely reproduced on the listing next to the accompanying instruction. It is usually used to explain the purpose or intention of a particular instruction or group of instructions.

Comment Line

An entire 80 column line may be utilized to print comments by coding an asterisk (*) in column 1. This entire line is merely reproduced on the assembly listing without further processing by the assembler.

SYMBOLS

Symbols are used in the label field of a symbolic machine instruction to identify that particular instruction and to represent its address. Symbols may be used for other purposes, such as the symbolic representation of some memory address, the symbolic representation of a constant, the symbolic representation of a register, etc.

No matter how the symbol is used, it must be defined. A symbol is defined when the assembler knows what value the symbol represents. There is only one way to define a symbol. The symbol must at some time appear either in the label field of an instruction or of an assembler directive. The symbol will be assigned the current value of the Location Counter when it appears in the label field of a machine instruction, or of an ORG, ACON, DATA, or RES pseudo-op. It may be assigned some other value through use of the EQU or SET assembler directives. A symbol may not appear in the name field more than once in a program, because this would cause the assembler to try to redefine an already defined label. The assembler will not do this and will flag all definitions of the particular label as an error. The only exception is that a symbol assigned a value by a SET pseudo-op may have another value assigned by subsequent SET pseudo-ops.

SYMBOLIC REFERENCES

Symbols may be used to refer to storage designations, register assignments, constants, etc. For example:

Address	Label	Operation	Operand
9B	MAZE	DATA	H'F5'
9C		LODA,3	MAZE

The symbolic label "MAZE" represents the address 9B. It is used in the machine instruction at address 9C to tell the assembler to build an instruction LODA,3 9B. The symbolic label, in this case, is a way for the programmer to specify an address without knowing exactly what the address should be when the program is written. In this example, assume there was a need to modify this sequence of code: a data statement was inserted between the original two statements.

Address	Label	Operation	Operand
99	MAZE	DATA	H'F5'
9A,9B		DATA	H'FE,3A'
9C		LODA,3	MAZE

Even though there was a program change which caused the data at MAZE to be located at address 99, the load instruction referencing the data did not have to be rewritten because the assembler could provide the proper physical address for the symbolic address MAZE. The instruction at address 9C will be assembled as LODA,3 99.

SYMBOLIC ADDRESSING

When writing instructions in the symbolic assembler language for the 2650, the addresses may be expressed through symbolic equivalents. The assembler will translate the symbolic address to its numeric equivalent during the assembly process.

It is good programming practice to make all address references symbolic, as this greatly eases the programmer's job in producing a working program. To make the register specification symbolic, one could equate a symbol to the register number:

```
RG3      EQU      3
          . . .
          . . .
          . . .
          LODA,RG3  MAZE
```

Forward References

A previously defined symbol is one which has appeared in the name field before it is referenced (as above). In contrast, a forward reference is a symbolic reference to a line of code when the symbol has not yet appeared in the name field. For example:

```
          ADDA,2    COEF
          . . .
          . . .
          . . .
COEF      DATA    D'123'
```

Forward references may be used anywhere in a program with the exception of the operand fields of the EQU, SET, RES, IF, and ORG statements.

Relative Addressing

The programmer may reference a memory cell either directly or via relative addressing. To refer directly to a memory cell of symbolic address MAIN, one has merely to use the name MAIN in the argument field of the referencing instruction. For example:

```
BIRA,R2  MAIN
```

It is also possible to express the address of a memory cell symbolically if some nearby cell is symbolically assigned. For example, to load the memory cell which is 5 cells higher in memory than the cell named MAIN, one need only to refer to it as MAIN+5:

```
LODA,2   MAIN+5
```

This latter method is called relative addressing. The relative count must be in the range -32,768 to +32,767.

The Location Counter and Symbol '\$'

There is one symbolic name, "\$," which is automatically defined by the assembler. This single character name is always symbolically equated to the contents of the assembler's location counter. Since the location counter is used by the assembler during the assembly process and is usually equated to the address of the byte being assembled, it represents the address of the instruction or data currently being specified. For example: BCTR,3 \$+5. The branch address will be interpreted by the assembler to be the address of the first byte of the branch instruction plus 5 bytes.

Hardware Relative Addressing

When using instructions which use "hardware relative addressing" (as distinguished from relative addressing discussed earlier in this section), it is important to realize the assembler will not only evaluate the expression which is given as an operand address, but will convert it to a hardware relative address (see the Hardware Specifications manual for a description of the addressing modes). For example:

Address	Name	Operation	Argument
100	SAM	LODA,R2	PAL
103		SUBI,R2	-3
105		BIRR,R3	SAM
107	next instruction		

In this code, the BIRR instruction specifies hardware relative addressing. Even though the equivalent value of the symbolic address SAM is 100, the relative addressing instruction requires a displacement relative to the address of the next sequential instruction. Therefore, the operand SAM will be evaluated as $-(\text{current location counter} + \text{length of BIRR instruction} - \text{SAM}) = -(105 + 2 - 100) = -(+7) = -7$. Remember, where the hardware instruction calls for "hardware relative addressing," the expression in the operand field will be evaluated as the displacement from the address of the next sequential instruction. The value of this displacement may range from -64 to +63.

Indirect Addressing

The symbol "*" is used to specify indirect addressing. For example:

	BCTA,3	*SAM
	...	
	...	
SAM	ACON	SUBR

In this sequence of instructions, the BCTA instruction specifies indirect addressing. The assembler will set the indirect bit (byte #1, bit #7) to '1' for this instruction.

Auto-Increment and Auto-Decrement

The symbols "+" and "-" are used to specify auto-increment and auto-decrement, respectively. For example:

```
LODA,RO    BUF,R3,+
```

In this instruction, which specifies auto-increment, the assembler sets bits #6 and #5 of byte #1 to "01." This option is specified in the instruction set tables as (,X).

IV. THE ASSEMBLY PROCESS

The 2650 assembler translates symbolic source code into machine language instructions. The assembler examines every source statement for syntactic validity and produces the equivalent machine code for the 2650 processor.

This is a two pass assembler, which means the entire source code is scanned twice by the assembler. On the first pass, all defined labels and their equivalent values are stored in a symbol table, the first byte of every instruction is fully determined, and some classes of errors are detected. During pass 2, symbolic address references are replaced by their values, all remaining errors are detected, and a listing and linkable object module are generated.

EXECUTION OF THE RELOCATABLE ASSEMBLER

The relocatable assembler is initiated by means of the SDOS command RASM. The format of the command is:

```
RASM,<infile>,[<list>], [<ofile>], [<print>]
```

Where:

- | | |
|-----------------------------|---|
| <code><infile></code> | denotes the file containing the assembler source text which is the user's source file if no macros are used. Otherwise, <code><infile></code> is the output file of the MACRO preprocessor. This is the only required parameter. |
| <code><list></code> | identifies the list output device. <code><list></code> may be a filename. If this parameter is omitted, no listing is produced. If CONO is specified, the listing lines are truncated after 79 characters. |
| <code><ofile></code> | identifies the file that is to contain the object program module. This file does not contain executable code. No object program module is stored if this parameter is omitted. |
| <code><print></code> | is a flag which determines whether or not macro expansions are to be printed. If <code><print> = G</code> , the expansions will be printed. The default is no print. A parameter other than G or (blank) produces an error message. |

Upon successful initiation, the relocatable assembler displays the message:

```
  ** 2650 RELOCATABLE ASSEMBLER VER n.n **
```

where:

n.n identifies the RASM version being executed.

LOCATION COUNTER

The assembler maintains a memory cell which it uses as a location counter. This location counter keeps track of the address of the next byte of storage to be allocated by the assembler. During coding, the programmer may think of the location counter as containing the address of the first byte of the instruction being written. In this assembler, the location counter is also used to provide load information.

ERROR DETECTION

During an assembly, the source program is checked for syntax errors. If errors are found, appropriate notification is given and the assembly proceeds. Although an assembled program containing errors generally will not run properly, it is considered good practice to complete the assembly to locate all errors at one time, rather than terminate it when an error is encountered.

ERROR CODES

There is a column on the listing in which an error indication may appear. Sometimes, because an error causes the assembler to view a subsequent statement incorrectly, a valid statement may be flagged as an error. A good rule is to fix errors in a particular line of code as they are discovered.

The following alphabetic characters are printed in the error indicator column and imply the corresponding message.

- A - Argument error. The argument has been coded in such a way that it cannot be resolved to a unique value.
- L - Label error. The label contains too many characters, contains invalid characters, has been previously defined, or is an invalid symbol.
- O - Op-code error. The op-code mnemonic has not been recognized as a valid mnemonic.
- P - Paging error. A memory access instruction has attempted to address across a page boundary.
- R - Register field error. The register field expression could not be evaluated, or when evaluated, was less than 0 or greater than 3, or the register field was not found.
- S - Syntax error. The instruction has violated some syntax rule.
- U - Undefined symbol. There is a symbol in the argument field which has not been previously defined.
- V - Numeric or expression stack overflow. Expression evaluation has resulted in an integer outside the valid range of -32,768 to +32,767, or the expression is too complex to evaluate. Processing of the source line is aborted.
- W - Warning. The assembler has detected a syntactically correct but unusual construction.

If more than one error occurs in a single statement, only the indicator for the first error encountered will be printed, and only the first error will be counted in the total errors.

I/O Errors

If an input/output error is detected, a message is displayed on the console and the assembly is terminated. The format of the message is:

```
** ERR <xx> <id>FILE **
```

where:

<xx> denotes the SRB status returned for the I/O operation.

<id> identifies the concerned file. The possible
 identifiers are:

 S - Source file, <ifile>

 L - Listing file, <list>

 O - Object file, <ofile>

Command Line Errors

If the optional parameter <print> is neither G nor a blank, the following message is displayed on CONO, and processing is aborted:

 INVALID 4TH PARAMETER

Symbol Table Error

A symbol table overflow produces the message:

 SYMBOL TABLE OVERFLOW

SYMBOL TABLE

An alphabetic symbol table is produced upon completion of RASM processing. This symbol table contains all program labels, their addresses, and attributes with which they are declared (i.e., as a label).

The maximum number of symbols allowed is 650 for a TWIN with 16K bytes of slave memory, and an additional 400 symbols for each 4K byte increment of memory.

The message SYMBOL TABLE OVERFLOW appears when the symbol table is full.

V. DIRECTIVES AND EXTENDED INSTRUCTION SET

ASSEMBLER DIRECTIVES

There are seventeen directives (pseudo-ops) which the assembler will recognize. These assembler directives, although written much like processor instructions, are simply commands to the assembler instead of to the processor. They direct the assembler to perform specific tasks during the assembly process, but have no meaning to the 2650 processor. These assembler directives are:

ORG
EQU
SET
ACON
DATA
RES
END
EJE
CEJE
PRT
SPC
TITL
PCH
REPRO
IF
ELSE
ENDIF

The use of the last three directives listed above (IF, ELSE, and ENDIF) is described in Section VI.

There are also three directives that relate to relocatability. These directives are:

CSECT
ENTRY
EXTRN

The use of these relocatable directives is described in Section VII.

ORG

Set Location Counter

The ORG directive sets the assembly time Location Counter to the location specified. The assembler assumes an ORG 0 at the beginning of the program if no ORG statement is given. The instruction ORG may not set the location counter to a value lower than the current value.

LABEL	OPERATION	OPERAND
{name}	ORG	expression

Where:

name optionally provides a symbol whose value will be equated to the value in the Location Counter after the ORG is evaluated.

expression when evaluated, results in a positive integer value. This value will replace the contents of the location counter, and bytes subsequently assembled will be assigned sequential memory addresses beginning with this value. Any symbols which appear in the argument must have been previously defined.

Examples:

```
LARR            ORG            YORD  
STAR           ORG           H'100'
```

EQU

Specify a Symbol Equivalence

The EQU directive tells the assembler to equate the symbol in the label field with the evaluable expression in the operand field. A symbol equivalence defined by an EQU directive cannot be redefined by a subsequent EQU or SET directive and cannot appear in the label field of a statement.

LABEL	OPERATION	OPERAND
name	EQU	expression

Where:

name is the symbol which is to be assigned some value by the execution of this directive.

expression is resolved to an integer value. If a symbol is used in the argument, it must have been previously defined.

Examples:

```
PAN      EQU      H'10F'  
LOP2     EQU      PAL  
RAMP     EQU      SLOP-3+PAL  
REG1     EQU      1
```

SET **Specify a Symbol Equivalence**

The SET directive tells the assembler to equate the symbol in the label field with the expression to be evaluated in the argument field.

The SET directive is identical to the EQU directive, except that the symbol defined by the SET directive may be redefined later in the program by another SET directive.

LABEL	OPERATION	OPERAND
name	SET	expression

The ACON directive tells the assembler to allocate two successive bytes of storage. The evaluated argument will be stored in the two bytes, the low-order 8 bits in the second byte and the high-order bits in the first byte. This directive is mainly intended to provide a double byte containing an address for use as the indirect address for any instruction executing in the indirect addressing mode. Any number of operands may be specified with one ACON directive. Each operand will be allocated two bytes of storage.

The data definitions in the instruction ACON may be preceded by a multiplication factor. This instruction cannot be used to define character strings.

NOTE

If an error is detected in the ACON instruction, two null bytes (00) are generated at the place where the error was detected. No further processing is done on the line. This is done to facilitate patching of the object code for debugging purposes.

LABEL	OPERATION	OPERAND
{ name }	ACON	expression

Where:

name is an optional label. If specified, the name becomes the symbolic address of the first byte allocated.

expression is some expression which must resolve to a positive value or zero. If positive, the value should be no larger than that which can be contained in two bytes. Otherwise, only the two least-significant bytes are used.

Example:

```
ASUB      ACON      SUBR
          ACON      SUBR,H'AFF0'
          ACON      3H'000F'
```

DATA

Defines Memory Data

The DATA directive tells the assembler to allocate the exact number of bytes required to hold the data specified in the argument field of this directive. Any number of bytes can be specified with one DATA directive.

The data arguments in the instruction DATA may be preceded by a multiplication factor. Additionally, the length of character strings can be defined explicitly.

Two types of data constants are distinguished: character strings and other constants. Both types may be specified in the same data instruction. A character string occupies a variable number of bytes; this number is either implicitly or explicitly specified. Other constants always occupy a single byte.

NOTE

If an error is detected in the DATA instruction, two null bytes (00) are generated at the place where the error was detected. No further processing is done on the line. This is done to facilitate patching of the object code for debugging purposes.

LABEL	OPERATION	OPERAND
{name}	DATA	argument

Where:

name is an optional label. If used, the name becomes the symbolic address of the first byte allocated by the directive.

argument is a general constant, a self-declining constant or a symbolic address. A multiple constant specification in the argument field will be handled as if the corresponding number of data statements were input.

To specify explicitly the length of the generated argument, precede the argument with the ASCII character L followed by a length value. Length value can be either an expression or a constant. The argument string will be truncated or blank-filled as necessary to satisfy the length value.

Examples:

```
PAL      DATA      LOOP,LOOP+1
         DATA      H'03,22,FC,A1'
         DATA      +127
         DATA      D'28'
DEFINE   DATA      A'THIS IS'
         DATA      AL10'THIS IS'
         DATA      AL2'THIS IS'
         DATA      3A'THIS IS'
         DATA      3AL2'THIS IS'
```

NOTE

If the argument evaluates to a value between 0 and 255, the result is an eight bit absolute binary number. DATA +127 results in H'7F'. Also, if the argument evaluates to a value which is less than 0, the result is a 2's complement, binary number. DATA H'-5' results in H'FB'. If the argument resolves to a value which requires more than one byte, only the least-significant byte is used, unless a length is specified.

RES **Reserve Memory Storage**

The RES directive tells the assembler to reserve contiguous bytes of storage. The number of bytes so reserved is determined by the argument. The reserved bytes are not set to a known value, but rather the effect of this directive is to increment the location counter.

LABEL	OPERATION	OPERAND
{name}	RES	expression

Where:

name is an optional label. If used, the name becomes the symbolic address of the first byte allocated.

expression is some evaluatable expression which must resolve to some positive integer or zero. If a symbol is specified, it must have been previously defined.

Example:

```
LOR            RES            23
MASK           RES           0'374'
                 RES           H'1A'
```

END

End of Assembly

The END directive informs the assembler that the last statement to be assembled has been input. The END directive causes the assembler to communicate the program start address to the object module. The assembler will not generate this instruction.

The operand of an END instruction, if present, must resolve to an address, and that address is used as the entry point of the module. Otherwise, the module has no explicitly defined entry point.

NOTE

A PHASE, as produced by the link editor needs an entry point and the entry point of the last module edited in a PHASE is used. When a module has no explicitly defined entry point, the begin address of the PHASE is taken as the entry point.

LABEL	OPERATION	OPERAND
	END	{expression}

Where:

expression is resolved to be the starting address of the program.

EJE

Eject the Listing Page

The EJE directive tells the assembler to advance the listing to the top of the next page regardless of the line position on the current listing page.

The directive is used primarily to organize listing for documentation purposes and does not appear in the listing.

LABEL	OPERATION	OPERAND
	EJE	

CEJE

Conditional Eject the Listing Page

This instruction is equivalent to a normal EJE instruction if the remaining number of lines on the current page is less than the value of the expression. If the remaining number of lines is equal to or greater than the value of the expression, the instruction has no effect.

LABEL	OPERATION	OPERAND
	CEJE	expression

Where:

expression

is an evaluable expression which must resolve to a positive integer.

PRT

Printer Control

The PRT directive tells the assembler to resume or discontinue printing of the assembled program.

This directive is used primarily to shorten assembly time by listing only that portion of the program which the user needs to see. This directive does not appear in the listing.

LABEL	OPERATION	OPERAND
	PRT	{ ON OFF GEN NOGEN }

Where:

- ON means print the assembled program
- OFF means discontinue print of the assembled program
- GEN means print generated macro expansions
- NOGEN means discontinue print of generated macro instruction expansions. When the G flag is entered in the RASM command line, it overrides all NOGEN flags.

NOTE

PRT is set to ON,NOGEN at the beginning of the assembly.

SPC

Space Control

The SPC directive tells the assembler to skip or space a number of lines.

This directive is used primarily to organize listings for documentation purposes and does not appear in the listing.

LABEL	OPERATION	OPERAND
	SPC	expression

Where:

expression

is an evaluable expression which must resolve to a positive integer. If the value of this expression is equal to, or greater than, the number of lines remaining on the page, the effect is the same as the EJE directive.

Example:

SPC 5

TITL

Title

The TITL directive tells the assembler to skip to the top of the next page and insert a given title into the main header.

This directive is used primarily for documentation purposes and does not appear in the listing.

LABEL	OPERATION	OPERAND
	TITL	expression

Where:

expression

is the title information not to exceed 38 character positions.

Example:

TITL MAIN PROGRAM

PCH

Punch Control

The PCH directive tells the assembler to selectively resume or discontinue the output of the object program module.

This directive is used primarily to shorten assembly time when an object program module is not desired or when only a portion of the object program module is desired.

LABEL	OPERATION	OPERAND
	PCH	{ ON } { OFF }

NOTE

PCH is set ON at the beginning of an assembly. When PCH OFF is specified, any prior object program module data is output.

REPRO

Copy Next Source Line

Often it is desirable to apply certain link editing commands to the link editor via the source file. Such commands are not automatically supplied by RASM. The REPRO directive is used to copy such commands from the source module. The instruction means copy the next source line into the object module and do not assemble that line. The instruction may not be labeled and has no operands.

The instruction REPRO is executed during pass 1 of an assembly run, so the reproduced statements appear at the beginning of the object module, before the external symbol dictionary (ESD).

LABEL	OPERATION	OPERAND
	REPRO	

Example:

```
REPRO  
INCLUDE FILEA
```

directs the link editor to include FILEA in front of this program object module in creating the current phase.

EXTENDED INSTRUCTION SET

In order to simplify the writing of programs and to increase the readability of programs, the mnemonic branch instruction op codes have been extended. The additional branch instructions are intended for use after compare, load and arithmetic instructions and the instruction TMI. These instructions with their standard equivalents are:

<u>Situation</u>	<u>Condition Code Value</u>	<u>Extended Instruction</u>	<u>2650 Equivalent Op Code</u>
After Compare	Branch if higher	BHR	BCTR,1
		BHA	BCTA,1
	equal	BER	BCTR,0
		BEA	BCTA,0
	lower	BLR	BCTR,2
		BLA	BCTA,2
	not higher	BNHR	BCFR,1
		BNHA	BCFA,1
	not equal	BNER	BCFR,0
		BNEA	BCFA,0
	not lower	BNLR	BCFR,2
		BNLA	BCFA,2
After Load and Arithmetic Operations	Branch if positive	BPR	BCTR,1
		BPA	BCTA,1
	zero	BZR	BCTR,0
		BZA	BCTA,0
	minus	BMR	BCTR,2
		BMA	BCTA,2
	not positive	BNPR	BCFR,1
		BNPA	BCFA,1
	not zero	BNZR	BCFR,0
		BNZA	BCFA,0
	not minus	BNMR	BCFR,2
		BNMA	BCFA,2
After TMI	Branch if ones	BOR	BCTR,0
		BOA	BCTA,0
	mixed	BMR	BCTR,2
		BMA	BCTA,2

Also, the following unconditional extended branch instructions have been added:

<u>Extended Instruction</u>	<u>2650 Equivalent Op Code</u>
BR	BCTR,3
BA	BCTA,3
BSR	BSTR,3
BSA	BSTA,3
RET	, RETC,3



VI. CONDITIONAL ASSEMBLY

The conditional assembly directives allow the programmer to vary the sequence of generated statements. Thus, the programmer can use these instructions to generate different sequences of statements from the same source program.

There are three conditional assembly directives. They are:

IF
ELSE
ENDIF

The format of the IF statement is:

LABEL	OPERATION	OPERAND
{name}	IF	expression

The standard expression rules apply with the addition of the following six relational operators which may be used in relational expressions:

.EQ. .NE. .LT. .GT. .LE. .GE.

The format for relational expression is

(expression.XX.expression)

Where:

.XX. is any one of the above relational operators.

The format of the ELSE statement is:

LABEL	OPERATION	OPERAND
	ELSE	

The format of the ENDIF statement is:

LABEL	OPERATION	OPERAND
	ENDIF	

Every IF statement must have a corresponding ENDIF statement. The use of the ELSE statement is optional, but if present, it must appear after the IF statement and before the ENDIF statement.

When an IF statement is encountered, the expression or relational expression is evaluated to be either true (not zero) or false (zero). If true, the following source statements are processed until an ENDIF is encountered. However, if an ELSE is encountered during this processing, the statements between the ELSE and the ENDIF are not processed. If false, the source statements following the IF are not processed until an ELSE or ENDIF is encountered, at which time normal processing resumes.

Conditional assembly constructs may be nested to eight levels but may not be overlapped. Therefore, the end of an inner IF construct (nested) must be encountered before the end of the outer IF construct is encountered.

The assembler listing will contain only those statements actually assembled. The statements between an IF and ENDIF (or ELSE) which do not produce object code do not appear on the listing.

If the optional label is included in the IF statement, the label will be assigned the value of the Location Counter when the next byte is actually assembled.

VII. RELOCATABILITY/RELOCATABLE ASSEMBLY

In order to allow links between separately assembled programs, each program object module begins with an external symbol dictionary (ESD). The ESD is printed at the beginning of the program listing. An ESD can contain the following types of records:

SD - Section definitions

The first record in each ESD is of the type SD. It contains the name of the section, its assembled origin and its length. The section name is specified by the CSECT instruction.

LD - Label definition

An ESD can contain several records of the type LD. They identify symbols declared as entry points into the program and their associated addresses. Other separately assembled programs may then refer to these symbols. The directive ENTRY is responsible for the generation of the LD records.

CD - Constant definitions

A CD record is similar to an LD record, except that the value is not an address but a constant. The CD records are also generated by the ENTRY directive.

ER - External reference

The ER records form the counterpart of the LD records and the CD records. They contain the symbols (with a sequence number) which are defined as external in the associated program. The directive EXTRN is used for this purpose.

At link editing time, a symbol table is constructed which contains all the symbols from the SD records, the LD records, and the CD records. Each symbol may occur only once. During the link editing of an object module, the values corresponding to its ER symbols are retrieved from this symbol table and edited into the object module.

The following RASM directives are available for the construction of an ESD.

CSECT

Identify Object Module

The instruction CSECT can be used to identify an object module. Each source program may contain at most one CSECT directive. Normally, this instruction is placed at the beginning of the program and format is:

LABEL	OPERATION	OPERAND
label	CSECT	

Where:

label

is optional; it simply assigns an identification to the object module for reference on the ESD. The address of the assembled origin of the module is assigned to the label. If there is no label, the ESD shows CSECT at the assembled origin of the object module without an ID. The directive has no operands.

ENTRY

Define Entry Points

The directive ENTRY is to be used to specify the symbols which may be referenced by other programs. The directive may not be labeled. The operand consists of a string of symbols separated by commas. The first blank encountered signals the end of the symbol string. These symbols must be previously defined as the label of some machine instruction or assembler instruction.

The format of the ENTRY statement is:

LABEL	OPERATION	OPERAND
	ENTRY	label list...

Where:

label list is one or more labels separated by commas.

Example:

```
ENTRY LABEL1  
ENTRY LABEL2, LABEL3
```

EXTRN**Define External Labels**

The directive EXTRN is used to specify the symbols which are defined in some other program but which are referred to in this program. EXTRN may not be labeled. The operand consists of a string of symbols separated by commas. The first blank encountered signals the end of the instruction.

These symbols may not be defined in the module containing the EXTRN directive. The only exception is given by EXTRN directives which are generated by means of a macro call expansion (see Section IX). Then the attribute external reference generated via the macro is ignored if the symbol is also defined elsewhere in the program.

The format of EXTRN statement is:

LABEL	OPERATION	OPERAND
	EXTRN	label list...

Where:

label list is one or more labels separated by commas.

Example:

```
EXTRN LABELA
EXTRN LABELB, LABELC
```

VIII. LINK EDITOR

The link editor, LINK, creates an executable load module from one or more object program modules generated by the relocatable assembler. Any object program module generated by RASM must be processed by LINK before it can be executed. The executable load module can be executed by the SDOS commands LOAD/GO or XEQ.

LINK is a two pass program that resolves inter-module references flagged by the relocatable assembler and converts relocatable addresses into absolute addresses. A load map is printed at the end of pass two that lists:

1. For each phase, the:

- Phase name (if any)
- Phase begin address
- Phase end address
- Phase entry point
- All object modules contained in that phase

2. For each object module, the:

- Module name (if such name exists)
- Module begin address
- All external symbols defined in the module with their value

3. Linkage editing error messages

When the link editor encounters an object program module ORGed at zero (either by default or by ORG 0) which will cross a 2K page boundary, it automatically relocates that module to the start of the next 2K page.

EXECUTION OF THE LINK EDITOR

The link editor is invoked by means of the SDOS command LINK. The format of the command is:

LINK,<input>, [<loadmap>], [<bfile>]

Where:

<input> contains commands to the link editor to locate the object program module or modules to be linked and identify the binary load module created. The input file may be the output from RASM containing link editor commands or a link editor command file.

<loadmap> specifies the device to which the load map is written. The default value for this parameter is no load map. <loadmap> may be a filename.

<bfile> is used to name the first binary load module when no PHASE command precedes the first INCLUDE command. When <bfile> is omitted and <input> contains no PHASE command, the edited output, i.e., the binary load module, is not stored.

After successful initiation, the program displays the message:

```
** LINK EDIT VER n.n **
```

Where:

n.n identifies the version of LINK being executed.

Commands to the link editor may be invoked as follows:

1. Include the commands in the RASM <ofile> file via the REPRO assembler directive in conjunction with LINK commands. Specify the RASM <ofile> as the LINK <input> .
2. Enter the LINK commands into a command file specified as <input> to LINK.
3. Combine 1 and 2 above, but the command file must be the one specified as LINK <input> .

ERROR MESSAGES

The link editor can display the following error messages: These error messages appear on the load map.

INVALID ASSEMBLED ORIGIN - The assembled origin of the program object module is an address less than the current location counter. The assembled origin is ignored and linking continues.

INVALID RECORD - An invalid command record is encountered. If the record is a null record (no characters in the line), processing aborts.

MULTIPLY DEFINED SYMBOL - An external symbol is defined in more than one object module

NESTING ERROR - Included modules are nested too deep. A nesting level up to 4 can be accepted. This means that the basic object program module may include object module (level 1). This object program module includes a next object program module (level 2), and so on to 4 levels.

PHASE COMMAND REQUIRED - The first two commands to LINK were INCLUDEs not preceded by a PHASE command. Processing is aborted

PAGING ERROR - An instruction with a 13-bit address refers to a location outside the current page.

TOO MANY PHASES - The object program file specified more than 20 phases.

UNDEFINED SYMBOL - A symbol is defined as EXTRN in one of the object modules, but it has not been defined as ENTRY in one of the other modules.

If an I/O error is detected, a message is displayed on CONO and processing aborts. The format of the message is:

** ERR <xx> <id>FILE **

Where:

<xx> denotes the SRB status return for the I/O operation

<id> identifies the concerned file. The possible identifiers are:

- I - Input file; LINK <input> , i.e., RASM <ofile> or LINK command file
- L - Load map file, <loadmap>
- B - Binary load module file, <bfile>

LINK EDITOR COMMANDS

Although the link editor is able to produce an executable load module from single object program modules, additional link editor commands can be useful. Except for the comment, all of the following link editor commands must be preceded by a blank.

PHASE

This command names a phase[†] to be created by the link editor. When two or more object programs are to be included in the phase, they must be preceded by a PHASE command. The format is:

PHASE<bfile>[,<origin>] [,<ID>]

Where:

_ (underscore) is a required space

<bfile> is the name of the binary phase module to be created.

<origin> specifies the begin address of the phase. It is represented by a normal expression. The phase address works in conjunction with the assembled origin, as follows:

[†]Phase used here is defined as separately stored load module. Such a module is equivalent to the modules produced by SDOS MODULE command.

<u>Phase</u>	<u>Assembled Origin</u>	<u>Resultant Address</u>
none	0	Current location counter (LC)
X	0	X
none	Y	Y
X	Y (X < Y)	a. Phase module begins at X b. CSECT begins at Y c. Arbitrary data fills $X \leq LC < Y$
X	Y (X > Y)	a. Phase module begins at X b. CSECT begins at X c. Warning message issued

<ID> is the module identification. It is a character string of, at most, 21 characters. Larger character strings are truncated without warning.

This command is used to include a separately assembled object program module. INCLUDE commands may be nested 4 levels. The format of the command is:

_INCLUDE<name>

Where:

_ (underscore) is a required space.

<name> is the name of the program module to be included. The link editor replaces this command with the program object module.

Comment

The input of the link editor may comprise comment lines, which are only listed in the load map. The symbol * in column one flags a line as a comment. The format of the comment is:

*** [comment]**

Note that the comment must not be preceded by a space.

All PHASE and INCLUDE commands are listed on CONO for each pass as they are processed by LINK.

IX. TWIN RESIDENT MACRO FACILITY

When programs are constructed, often the same or nearly the same sequences of instructions are written. If this occurs, macro instructions can be used to reduce the number of written source code instructions. When a macro call is processed, it is replaced by a sequence of machine instructions from a corresponding macro definition. The macro definition contains the data to be generated.

A single level, or non-nested, macro processing facility is available with the relocatable assembler. It is implemented as a preprocessor; that is, from the source program, it produces a text file for use as the source input to the relocatable assembler, RASM.

EXECUTION OF THE MACRO PROCESSOR

The macro pre-processor is invoked with the SDOS command MAC, as follows:

MAC,<source>,<expanded source>, [<macro file>], [<error list>]

Where:

- | | |
|-------------------|--|
| <source> | is the user program file containing the macro calls. |
| <expanded source> | is the output file of the macro preprocessor to be used as the input source file to RASM. |
| <macro file> | is the file containing user written macro definitions for inclusion in the source program. |
| <error list> | is the optional list device or filename to which error messages are listed. The default is CONO. |

Macro Libraries

Two macro libraries may contain input to the macro processor:

1. <macro file> is the user's macro definition file and, when specified in the command, is searched first.
2. SYSMACRO is a default macro library. SYSMACRO is searched whenever a macro definition is not found in <macro file> or if a <macro file> is not specified. SYSMACRO must reside on disk drive 0.

Macro Processing

The macro preprocessor displays the following sign-on message:

```
** MACRO PREPROCESSOR VER n.n **
```

Where:

n.n identifies the version of MAC being executed.

Then, the input files are opened and macro definitions from both macro definition files are read into a memory buffer. If no errors have been detected at this point, the <source> file is processed.

When an end-of-file (EOF) condition is detected in <source> , processing is terminated.

ERROR MESSAGES

Macro preprocessor error messages are written to CONO, or to the specified <error list> file.

Certain error conditions are detected in the macro definition file before the source file is processed. Checking is done on macro names and on the size of the macro definition files. Error messages issued for these error conditions are:

INVALID MACRO NAME

The macro name is more than five characters or the macro prototype contains no macro name.

TOO MANY MACRO DEFINITIONS

The internal macro storage table or macro name table has overflowed. Up to 64 macro names are allowed. This error causes processing to abort.

NO MACRO DEFS

No user macro definition library was specified, and SYSMACRO was not on the disk drive 0. Processing aborts.

When an error has been detected during the processing of a macro instruction or control instruction, an alphabetic character indicating error type is printed with the source file line and corresponding macro definition line at which the error occurred. The format of this two line message is:

```
<C><source macro call>
```

```
<macro definition line>
```

Where:

<C> is the error indicator, as follows:

- C - Invalid control instruction
- P - Error in prototype
- D - Error in macro definition
- I - Error in macro instruction
- O - Output line becomes too long

<source macro call> represents the concerned source file input line

If an input/output error occurs, a message is displayed and the program is terminated. The format of the message is:

```
** ERR <xx> <id>FILE **
```

Where:

<xx> represents the SDOS SRB status returned for the I/O operation.

<id> represents the concerned file. Possible values of ID are:

I - Input file, <source>

O - Output file, <expanded source>

M - Macro definition file, <macro file> or SYSMACRO

<source> FILE CONTENTS

The <source> file can contain the following type of lines:

1. Control instructions. The first character of a control instruction is a period. The only recognized control instruction in the <source> file is:

```
.TABS N1 [,N2,N3...N7] [ comment ]
```

Where:

N_n is a decimal number in the range 0 - 70 and specifies a tab stop for the output lines. At most, 7 tab stops are accepted. The default tab stops are 10, 20, 30, 40, 50, 60 and 70. To reset tab stops to the default values, .TABS may be entered without parameters.

2. Comment lines. The first character of a comment line is *. These lines are added to the output file without any processing.
3. Data lines. The contents of the opcode field of assembly language source code are compared with the <macroname> of all macro definition prototypes. Upon a match, the corresponding macro definition is interpreted. Otherwise, the data line is simply copied to the output file.

<macro file> CONTENTS

The <macro file> consists of commands to the macro preprocessor and macro definitions.

MACRO DEFINITIONS

Macro definitions are bracketed by the macro preprocessor control instructions .MACRO and .MEND. The commands must be the only text on the line, and the period must appear in column one. The macro definition consists of one prototype statement followed by the macro body.

Prototype

The prototype defines the macro name and the macro definition formal parameters. It must be the first statement of a macro definition following the .MACRO control instruction and has the form:

[&<label>] <macroname> [&<param>, &<param>, ...]

The <macroname> can be at most five (5) characters long and may not begin in column 1.

Macro Body

The macro body follows the prototype and consists of assembly language statements and macro control instructions. Nesting of macros is not supported. Where formal parameters declared on the prototype line appear in the macro body, they are replaced by the actual parameter value passed in the macro call.

Macro control instructions do not appear in the macro expansion, but give directions to the macro processor as to how to do the expansion. All macro control instructions begin with a period and may be labeled with macro labels.

MACRO VARIABLES

During the processing of macro calls, the text corresponding to the macro body is generated. Macro labels, macro control instructions, and the prototype will not appear in the expanded output. Also, expressions and variables are replaced by their actual computed values.

Identifiers of up to six characters can be used as macro variables. Two types of variables are recognized:

- Binary variables can assume numeric values from -32,768 to +32,767.
- Character variables may be a string of up to 14 alphanumeric characters. Empty strings are valid.

These variables can be declared explicitly by the DECLARE macro command or implicitly by their appearance in a Prototype statement.

A leading & identifies macro variables.

Parameters are simply variables that are implicitly declared by their appearance in the prototype statement.

All special characters, except `_` (underscore) are treated as normal data characters. The character `_` (underscore) at the end of a variable is used as concatenation operator; it will not appear in the output line.

EXPRESSIONS

Expressions may occur in macro control instructions as well as in assembly language statements. Expressions are nearly identical to those used by RASM, but the operators `<` (less than) and `>` (greater than) do not exist. An expression is constructed with operators, parentheses and factors.

Constant

The following types of constants are possible:

- H - Hexadecimal constants
- O - Octal constants
- B - Binary constants
- D - Decimal constants
- Self defining Constants

All constants, except self defining constants, are enclosed by single quotation marks and are preceded by the appropriate code character, as defined previously in this manual.

A constant is internally expressed as 15 bits plus a sign bit and it can have a value in the range -32,768 through +32,767.

Character String

A character string is enclosed by quotation marks and is preceded by the code character A. A character string may contain up to 14 characters. Any character, except a single quotation mark, may occur within a character string. A quotation mark in a character string is to be specified by two consecutive single quotation marks. These two characters are then interpreted as a single character and not as a string delimiter.

Variable

A variable is identified by & as the first character. The value of a variable is a character string or a constant.

Expression Between Parentheses

The operators, ordered on increasing priority, are:

.OR. .XOR.	Logical OR and Exclusive OR
.AND.	Logical AND
.NOT.	Logical NOT (1-Complement)
.EQ. .NE. .GE. .LE. .LT.	Relations
+ -	Arithmetic addition and subtraction
* / .SHL. .SHR.	Arithmetic multiplication, division and shifts

The requirements and results of the various operations are:

- Arithmetic operations - The operands are binary numbers or it must be possible to convert the operands into binary numbers. The result is a binary number which shall be in the range or -32,768 to +32,767.
- Relations - The operands are character strings or binary numbers. Negative numbers are preceded by a minus sign, and leading zeros are omitted. These character strings are compared from left to right and the result of the comparison is:

1. Equal

Both strings have the same length and contain the same characters in an identical sequence.

2. Greater Than

The number of characters of the second operand is less than the number of characters of the first operand or the binary number, which represents the current character of the first operand, is greater than the binary number which represents the current character of the second operand.

3. Less Than

The length of the second operand exceeds the length of the first operand or the binary number, corresponding to the current character of the first operand, is less than the binary number corresponding to the current character of the second operand.

The final result of a relation operation is a binary number conforming to the following matrix:

	EQ	NE	LE	LT	GE	GT
Equal	1	0	1	0	1	0
Greater Than	0	1	0	0	1	1
Less Than	0	1	1	1	0	0

Logical Operations

The operands of logical operations are binary numbers. Character strings are, if applicable, converted to binary numbers.

MACRO PARAMETER EXPANSION

A macro call is simply a reference to a macro definition. It may contain the actual values of any formal parameters defined in the macro prototype statement.

In the macro definition a leading & identifies formal parameters. Three types of parameters are distinguished:

1. Label parameters. A label may be attached to the macro expansion by use of the label parameter. The formal parameter &<label> in the prototype statement passes the label on the macro call to the macro expansion. The label parameter must begin in column one of the macro prototype, and is of the form:

&<label>

2. Positional parameters. The values of positional parameters are given by the character strings at the associated positions in the macro prototype. Positional parameters appear in the operand field and have the form:

&<param>

3. Keyword parameters. The actual values of keyword parameters are identified by the corresponding keywords in the macro calls and are declared in the operand field of the macro prototype. Keyword parameters do not have pre-defined positions, and have the form:

&<param> = <value>

Where:

<value> is the parameter default value when the parameter is not specified in the macro call. <value> may be blank.

When both keyword and positional parameters are used in one macro call, the positional parameters are to be listed first.

Example:

The following macro moves a string of specified length.

```
.MACRO                                Begin macro
&L      MOVE   &A,&B,&N                Prototype statement
        LODI,R1 &N
&L      SUBI,R1 1
        LODA,RO &A,R1
        STRA,RO, &B,R1
        BRNR,R1 &L
        }                               Macro body
.MEND                                     End macro
```

The corresponding macro call contains a label parameter and the values of three positional parameters:

```
        LBL      MOVE      P,Q,3
```

The character string following the first blank in the call, MOVE, identifies the <macroname> . It corresponds to the operation code of a normal machine instruction. The character strings LBL, P, Q and 3 represent the actual values of formal parameters in the associated macro definition for MOVE.

Processing of the above macro call results in the following output:

```
        LODI,R1 3
LBL     SUBI,R1 1
        LODA,RO P,R1
        STRA,RO Q,R1
        BRNR,R1 LBL
```

Note that the control instructions and the prototype do not appear in the output. The formal parameters have been replaced by their actual values in the expanded macro call.

Suppose most of the string moves in our program move one byte of data to Q. The macro may then be written:

```
.MACRO
&L      MOVE      &A,&B=Q,&N=1
        LODI,R1   &N
&L      SUBI,R1   1
        LODA,RO   &A,R1
        STRA,RO   &B,R1
        BRNR,R1  &L
.MEND
```

The macro call now need only specify the data source as follows:

```
                LBL      MOVE      C
```

resulting in this macro expansion:

```
                LODI,R1   1
LBL            SUBI,R1   1
                LODA,RO   C,R1
                STRA,RO   Q,R1
                BRNR,R1  LBL.
```

Note that in this second example the keywords for string length and destination are N and B respectively. To move a string of length 3, simply specify the "length" keyword and actual value in the macro call.

```
                LBL      MOVE      P,N=3
```

The expansion of the above macro call results in code identical to that shown in the first example.

MACRO COMMANDS

Macro definitions can contain special instructions to the macro processor to allow greater flexibility at expansion time. Macro commands begin with a period and those whose first character is underscored in this manual can be abbreviated to the first letter of the command. Macro commands must appear only one per line of text.

.&<macrolabel>

This label, within the macro definition, labels macro commands. Note that this is not the same as the formal <label> parameter.

.ASSIGN &<variable> = <expression>

A simple assignment statement. The value of the expression is assigned to the variable.

.DECLARE &<variable>,...

This declares variables for use by the macro processor. Upon declaration, the variables are initialized with the empty character string.

**.GENERATE &<variable> = <expression1>, <expression2>
.END**

The GENERATE - END pair delimits a group of statements to be processed iteratively. The GENERATE statement declares a new variable, which assumes the value of the first expression as its initial value. The variable is incremented by one each time the END statement is encountered. When the value of the variable exceeds that of the second expression, the iteration stops. Variables declared with a GENERATE loop exist only inside that loop.

.IF <expression> .&<macrolabel>

This command conditionally includes a block of text into the macro expansion. If <expression> evaluates to true, the next statement to be considered for inclusion is at .&<macrolabel> . If <expression> is false, the next statement is the one following the IF statement. Representations of true and false are as follows: binary: 0 = false
not 0 = true

character: all blanks = false
at least 1 non-blank in the string = true

.JUMP .&<macrolabel>

The next statement to be scanned for inclusion is the one indicated by .&<macrolabel> .

.MACRO
⋮
.MEND } macro definition

This pair of commands must bracket each macro definition and cannot be abbreviated. They must begin in column 1; that is, **.MEND** cannot be labeled.

.MESSAGE <text>

This command causes <text> to be displayed on the system console, CONO.

.NOP

No operation. This is used for defining macrolabels, since they cannot appear on a blank line, or for jumping to the end of a macro definition since **.MEND** cannot be labeled. For example,

```
.J.&LBL  
.  
.  
.&LBL .NOP
```

.SUBSTRING &<variable> = &<variable>(<expression1>,<expression2>)

A substring of the second variable is assigned to the first variable. The first expression indicates the beginning character of the substring (starting with the value zero for the first character of the string), while the second expression gives the length of the substring.

.TAB <character>

This sets the tab character. The definition remains in effect only during the current macro definition. & is not a valid tab character.

*<text>

This is a comment; the line is ignored.

00

0

0

0

00

X. SAMPLE PROGRAM SET

Included here are two 2650 assembly language source programs, a SYSMACRO file listing, a LINK command file listing and the loadmap resulting from linking the assembled program object code from the two source programs.

The sample programs consist of a main program and a subroutine for an arithmetic bubble sort. These programs are called MAIN;S and SUBR;S respectively. The SYSMACRO file contains two macro definitions:

1. Register value equates
2. Condition code value equates

To expand the macro calls, REGS and CODES, for the equate values run:

```
MAC MAIN;S MAIN;E
MAC SUBR;S SUBR;E
```

Since no user macro library is specified in the MAC command lines, MAC will automatically look in SYSMACRO for macro definitions. The expanded source file names are denoted by the ";E" suffix on the filenames.

Now, the source files with expanded macros can be assembled. To do this, enter:

```
RASM MAIN;E LPT1 MAIN
RASM SUBR;E LPT1 SUBR
```

We have directed the assembly list output to the line printer and the assembled object program code to files MAIN and SUBR, both of which may now be input to the link editor in order to create an executable binary load module. To link these object program files, enter:

```
LINK LNKSRT LPT1
```

LNKSRT is a small file containing only commands for the link editor. Note that the PHASE command specifies the name of the binary load module as BUBBLE and it begins at address H'0000'. The included main program, MAIN, itself contains an INCLUDE command for the subroutine object program file. This causes the subroutine object program file to be linked first, in front of the MAIN program. The module entry address, specified by the main program END statement, is resolved to H'00D8'. Refer to the loadmap listing for particulars regarding resolution of addresses by the link editor.

The binary load module, BUBBLE, may now be loaded and executed by entering the SDOS commands XEQ or LOAD and GO.

ASSEMBLY LISTING

Figures 10-1 through 10-5 are sample program listings produced by RASM. The following explanations are keyed to the listing.

1. Page heading - Displays the current versions and level of RASM and the title, if any, specified by the TITL assembler directive.
2. Page number - Every page of the listing is numbered sequentially.
3. Location column - The numbers in this column are equal to the value of the assembly Location Counter and indicate the address at which the first byte is located.
4. Object - This field describes the data bytes which are stored sequentially starting at the address in the Address Column.
5. Address column - The numbers in this column are the results of evaluating address expressions which are in the instruction operand field. For the EQU and SET directives, this column contains the hexadecimal value of the operand field.
6. Error column - This column may contain an error code as detailed elsewhere in this chapter.
7. Statement Number - This number corresponds to the line number of the source program file.
8. Source code - This area of the listing reproduces the source code as it was read by the assembler.
9. Total errors - This field indicates the total number of errors detected by the assembler during passes 1 and 2 of the assembly process.

```
1 *
2 * SYSMACRO FILE - CONTAINS STANDARD 2650 REGISTER AND CONDITION CODE EQUATES
3 *
4 *
5 * REGISTER EQUATES
6 *
7 . MACRO
8 REGS
9 . T /
10 R0/EQU/0
11 R1/EQU/1
12 R2/EQU/2
13 R3/EQU/3
14 . MEND
15 *
16 * CONDITION CODE EQUATES
17 *
18 . MACRO
19 CODES
20 . T /
21 EQ/EQU/0
22 GT/EQU/1
23 LT/EQU/2
24 Z/EQU/0
25 P/EQU/1
26 N/EQU/2
27 UN/EQU/3
28 . MEND
29 *
30 * END OF SYSMACRO FILE
31 *
```

Figure 10-1. SYSMACRO File

TWIN 2650 ASSEM VER 3.0 EXTERNAL SYMBOL DICTIONARY

SYMBOL	TYPE	VALUE	LENGTH
BMAIN	SD	0000	0013
CNT	ER	0002	
SORT	LD	0005	
SUB1	ER	0001	

Figure 10-2. MAIN Program Assembly (Page 1)

LOC	OBJECT	ADDR	E	STMT	SOURCE	LINE
				2	BMAIN	CSECT
				3	*	
				4	* ARITHMATIC BUBBLE SORT PROGRAM	
				5	*	
				6	* GET REGISTER AND CONDITION CODE EQUATES FROM SYSMACRO.	
				7	PRT	NOGEN THE "G" PARAM WILL OVERRIDE THIS FLAG.
				8	REGS	
				13	CODES	
				21	PRT	GEN
				22	*	
				23	* MAIN PROGRAM	
				24	*	
0000	0F0012	0012		25	STRT	LODA, R3 LEN LOAD BUFFER LENGTH IN R3
0003	7502			26	CPSL	2 SET ARITHMATIC COMPARISONS
0005	A701			27	SORT	SUBI, R3 1 DECREMENT LOOP COUNTER
0007	CF0000	0000		28	STRA, R3	CNT STORE LOOP COUNTER
000A	1A05	0011		29	BCTR, N	STOP IF COUNT DECREMENT PAST 0. STOP
000C	3F0000	0000		30	BSA	SUB1
000F	1B74	0005		31	BR	SORT
0011	40			32	STOP	HALT
				33	*	
0012	0F			34	LEN	DATA 15 LENGTH OF BUFFER
				35	*	
				36	REPRO	COMMAND TO LINK EDITOR
				37	INCLUDE	SUBR
				38	ENTRY	SORT
				39	EXTRN	SUB1, CNT
		0000		40	END	STRT

Figure 10-2. MAIN Program Assembly (Page 2)

SYMBOL	TYPE	VALUE	LENGTH
BMAIN	SD	0000	0013
BMAIN	LD	0000	
CNT	ER	0002	
E0	CD	0000	
GT	CD	0001	
LT	CD	0002	
LEN	LD	0012	
N	CD	0002	
P	CD	0001	
R0	CD	0000	
R1	CD	0001	
R2	CD	0002	
R3	CD	0003	
STRT	LD	0000	
SORT	LD	0005	
STOP	LD	0011	
SUB1	ER	0001	
UN	CD	0003	
Z	CD	0000	

TOTAL NUMBER OF ERRORS 0

Figure 10-2. MAIN Program Assembly (Page 3)

SYMBOL	TYPE	VALUE	LENGTH
BSUBR	SD	0000	0008
CNT	LD	00C8	
LOOP	LD	0004	
SUB1	LD	0001	

Figure 10-3. SUBR Program Assembly (Page 1)

LOC	OBJECT	ADDR	E	STMT	SOURCE	LINE
					1	* SUBROUTINE FOR BUBBLE SORT.
					2	*
					3	*
					4	* GET REGISTER AND CONDITION CODE EQUATES FROM SYSMACRO.
					5	*
					6	REGS
					11	CODES
					20	B SUBR CSECT
0000	00				21	ZERO DATA 0
					22	*
					23	* DOES ONE ITERATION THROUGH BUFFER
					24	*
0001	0E0000	0000			25	SUB1 LODA, R2 ZERO R2 COUNTS COMPARISONS
0004	EE00C8	00C8			26	LOOP COMA, R2 CNT IF =, ITERATION COMPLETE
0007	14				27	RETC, EQ
0008	0E60C9	00C9			28	LODA, R0 BUF, R2 LOAD FIRST # OF THE PAIR
000B	EE20C9	00C9			29	COMA, R0 BUF, R2, + COMPARE WITH SECOND #
000E	9974	0004			30	BCFR, GT LOOP IF FIRST LT OR =, LOOP BACK
0010	C1				31	STRZ R1 MOVE LARGER # TO R1
0011	0E60C9	00C9			32	LODA, R0 BUF, R2 LOAD SMALLER # INTO R0
0014	CE60C8	00C8			33	STRA, R0 BUF-1, R2 STORE SMALLER # IN FIRST
0017	01				34	LODZ R1 MOVE LARGER # TO R0
0018	CE60C9	00C9			35	STRA, R0 BUF, R2 STORE LARGER # SECOND
001B	1B67	0004			36	BCTR, UN LOOP LOOP BACK
					37	*
					38	*
00C8					39	ORG 200
00C8					40	CNT RES 1
00C9					41	BUF RES 15 BUFFER TO BE SORTED
					42	*
					43	ENTRY SUB1, LOOP, CNT
					44	END

Figure 10-3. SUBR Program Assembly (Page 2)

SYMBOL	TYPE	VALUE	LENGTH
BSUBR	SD	0000	0008
BSUBR	LD	0000	
BUF	LD	00C9	
CNT	LD	00C8	
EQ	CD	0000	
GT	CD	0001	
LT	CD	0002	
LOOP	LD	0004	
N	CD	0002	
P	CD	0001	
R0	CD	0000	
R1	CD	0001	
R2	CD	0002	
R3	CD	0003	
SUB1	LD	0001	
UN	CD	0003	
Z	CD	0000	
ZERO	LD	0000	

TOTAL NUMBER OF ERRORS 0

Figure 10-3. SUBR Program Assembly (Page 3)

TWIN 2650 LINK EDITOR VER 3.0 LOAD MAP

PHASE	BEGIN	END	ENTRY	SYMBOL	TYPE	VALUE	OFFSET	PAGE	1
-------	-------	-----	-------	--------	------	-------	--------	------	---

*
* LNKSRT - LINK MAIN PROGRAM AND SUBROUTINE FOR BUBBLE SORT
*

BUBBLE	0000	00EA	00D8	BSUBR	CSECT	0000	0000		
				CNT	ENTRY	00C8			
				* LOOP	ENTRY	0004			
				SUB1	ENTRY	0001			
				BMAIN	CSECT	00D8	00D8		
				* SORT	ENTRY	00D0			

Figure 10-4. Load Map

```
1 *  
2 * LNKSR - LINK MAIN PROGRAM AND SUBROUTINE FOR BUBBLE SORT  
3 *  
4 PHASE BUBBLE, H'0000', BUBBLE SORT PROGRAM  
5 INCLUDE MAIN
```

Figure 10-5. LINK Command File

APPENDIX A MESSAGES

Error Message	Meaning	Recovery Procedure
<u>ASSEMBLY</u>		
A - Argument error	The argument has been coded in such a way that it cannot be resolved to a unique value.	Self-explanatory
L - Label error	The label contains too many characters, contains invalid characters, has been previously defined, or is an invalid symbol.	Self-explanatory
O - Op-code error	The op-code mnemonic has not been recognized as a valid mnemonic.	Self-explanatory
P - Paging error	A memory access instruction has attempted to address across a page boundary.	Self-explanatory
R - Register field error	The register field expression could not be evaluated, or when evaluated, was less than 0 or greater than 3, or the register field was not found.	Self-explanatory
S - Syntax error	The instruction has violated some syntax rule.	Self-explanatory
U - Undefined symbol	There is a symbol in the argument field which has not been previously defined.	Self-explanatory
V - Overflow	Numeric or expression stack overflow.	Self-explanatory
W - Warning	The assembler has detected a syntactically correct but unusual construction.	Self-explanatory

Error Message	Meaning	Recovery Procedure
<u>INPUT/OUTPUT</u>		
ERR<xx><ID>FILE	<p>And I/O error has occurred and the program is terminated.</p> <p><xx> denoted for SRB status returned for the I/O operation.</p> <p><ID> identifies the concerned file. The possible identifiers are:</p> <ul style="list-style-type: none"> S - Source file L - Listing file O - Object program file 	<p>Correct the problem in the named file according to the SRB status code (the list is contained in this appendix).</p>
INVALID 4th PARAMETER	<p>Optional macro expansion print parameter is neither G nor blank. Processing is aborted.</p>	<p>Reenter the RASM command with the correct parameter.</p>
SYMBOL TABLE OVERFLOW	<p>Symbol table is full.</p>	<p>Reduce the number of symbols in the program source file and reassemble.</p>
<u>LINK EDITOR</u>		
NO PHASE	<p>When the first two link commands are INCLUDE, a PHASE is required preceding them.</p>	<p>Insert a PHASE card in front of the LINK command stream and reLINK the files.</p>
INVALID RECORD	<p>An invalid command record is encountered.</p>	<p>Self-explanatory</p>
MULTIPLY DEFINED SYMBOL	<p>An external symbol is defined in more than one object program module.</p>	<p>Change the symbol name in one of the object program modules.</p>
NESTING ERROR	<p>Included modules are too deep. A nesting level up to 4 can be accepted.</p>	<p>Adjust INCLUDE nesting to 4 levels or less.</p>
PAGING ERROR	<p>An instruction with a 13-bit address refers to a location outside the current page.</p>	<p>Self-explanatory</p>

Error Message	Meaning	Recovery Procedure
<u>LINK EDITOR (Cont'd)</u>		
TOO MANY PHASES	The object file specifies more than 20 phases.	Link to intermediate phases, then link these to the final load module.
UNDEFINED SYMBOL	A symbol is defined as EXTRN in one of the object modules, but it has not been defined as ENTRY in one of the other modules.	Non-fatal error.
ERR <xx> <ID> FILE	An I/O error has occurred and the program is terminated. <xx> represents the SDOS SRB status. <ID> represents the concerned file. Possible values of <ID> are: I - One of the input files L - Load map file/device B - Binary load module output	Correct the problem in the named file according to the SRB status (the list is contained in this appendix).
<u>MACRO FACILITY</u>		
<C> <source macro call> <macro definition>	An error has been detected during the processing of an instruction. where <C> : C - Invalid control instruction P - Error in prototype D - Error in macro definition I - Error in macro instruction O - Output line becomes too long	Correct the line in error as dictated by the <C> error flag.

Error Message	Meaning	Recovery Procedure
<u>MACRO FACILITY (Cont'd)</u>		
ERR <xx> <ID> FILE	<p>An input/output error has occurred and the program is terminated.</p> <p><xx> represent the SDOS SRB status</p> <p><ID> represents the concerned file.</p> <p>Possible values of <ID> are:</p> <ul style="list-style-type: none"> I - Input source file O - Output expanded source file M - One of the macro definition files 	<p>Correct the problem in the named file to the SRB status code (the list is contained in this appendix).</p>
INVALID MACRO NAME	<p>The macro name is more than five characters or the macro prototype contains no macro name.</p>	<p>Shorten the macro name to five characters or enter a valid prototype statement.</p>
TOO MANY MACRO DEFINITIONS	<p>The internal macro storage table or macro name table has overflowed. Up to 64 macro names are allowed. This error causes processing to abort.</p>	<p>Split the macro definition file and run the source file against each smaller macro definition file.</p>
NO MACRO DEFS	<p>No user library was specified and SYSMACRO was not found on drive 0. Processing is aborted.</p>	<p>Rerun MAC. Specify the correct macro library file.</p>

SDOS SRB Status Codes

Status	Description
0	Function complete/no error
1	New file
2	Illegal channel number
3	Channel not assigned
4	Channel busy
5	Illegal function code
6	Short read
7	Short write
8	Illegal drive number
9	File in use
A	Device not operational
B	Device not available
C	Device not ready
D	Device in use
E	Directory read error
F	Directory write error
10	Directory full
11	Device read error
12	Device write error
13	Invalid address: Utility routine conflict or memory wrap-around
14	Unused

SDOS SRB Status Codes (Cont'd)

Status	Description
15	File name in use
16	Illegal file name
17	File in R/W progress
18	Channel already assigned
19	Incorrect diskette
7F	I/O in progress
FF	End of file/end of device

APPENDIX B

SUMMARY OF 2650 INSTRUCTION MNEMONICS

In these tables parentheses are used to indicate options. In no case are they coded in any instruction. The following abbreviations are used:

- r— register expression, must evaluate to $0 \leq r \leq 3$.
- v— value expression
- *— indirect indicator
- a— address expression
- x— index register expression
- X— index register expression with optional auto-increment or auto-decrement

Note

- The use of the indirect indicator is always optional.
- When an index register expression is specified, it can be followed by '+', '-' which indicates use of auto-increment or auto-decrement of the index register. Example:

LODA,0 DPR,R3,+

BXA, BSXA are exceptions and do not permit auto-increment or auto-decrement.

- Even though an address expression is specified in a hardware relative addressing instruction, the assembler develops it into a value of $(-64 \leq v \leq +63)$.
- A memory reference instruction which requires indexing may use only register 0 as the destination of the operation.
- If an index register expression is used with either BXA or BSXA instructions, it must specify index register #3 (either register bank) for indexing. Any other value in the index field will produce an error during assembly. However, it is not necessary to use an index register expression with these instructions: a blank in this field will default to register 3.

LOAD/STORE INSTRUCTIONS			Length (bytes)	SUBROUTINE BRANCH/RETURN INSTRUCTIONS			Length (bytes)
LODZ	r	Load Register Zero	1	BSTR,v	(*a)	Branch to Subroutine on Condition True, Relative	2
LODI,r	v	Load Immediate	2	BSFR,v	(*a)	Branch to Subroutine on Condition False, Relative	2
LODR,r	(*a)	Load Relative	2	BSTA,v	(*a)	Branch to Subroutine on Condition True, Absolute	3
LODA,r	(*a),(X)	Load Absolute	3	BSFA,v	(*a)	Branch to Subroutine on Condition False, Absolute	3
STRZ	r	Store Register Zero	1	BSNR,r	(*a)	Branch to Subroutine on Non-Zero Register, Relative	2
STRR,r	(*a)	Store Relative	2	BSNA,r	(*a)	Branch to Subroutine on Non-Zero Register, Absolute	3
STRA,r	(*a),(X)	Store Absolute	3	BSXA	(*a),(x)	Branch to Subroutine, Indexed, Unconditional	3
ARITHMETIC INSTRUCTIONS				RETC,v		Return From Subroutine, Conditional	1
ADDZ	r	Add to Register Zero	1	RETE,v		Return From Subroutine and Enable Interrupt, Conditional	1
ADDI,r	v	Add Immediate	2	ZBSR	(*),a	Zero Branch to Subroutine Relative, Unconditional	2
ADDR,r	(*a)	Add Relative	2	PROGRAM STATUS INSTRUCTIONS			
ADDA,r	(*a),(X)	Add Absolute	3	LPSU		Load Program Status, Upper	1
SUBZ	r	Subtract from Register Zero	1	LPSL		Load Program Status, Lower	1
SUBI,r	v	Subtract Immediate	2	SPSU		Store Program Status, Upper	1
SUBR,r	(*a)	Subtract Relative	2	SPSL		Store Program Status, Lower	1
SUBA,r	(*a),(X)	Subtract Absolute	3	LDPL ¹		Load Program Status, Lower	3
LOGICAL INSTRUCTIONS				STPL ¹		Store Program Status, Lower	3
ANDZ	r	And to Register Zero	1	CPSU	v	Clear Program Status, Upper, Selective	2
ANDI,r	v	And Immediate	2	CPSL	v	Clear Program Status, Lower, Selective	2
ANDR,r	(*a)	And Relative	2	PPSU	v	Preset Program Status, Upper, Selective	2
ANDA,r	(*a),(X)	And Absolute	3	PPSL	v	Preset Program Status, Lower, Selective	2
IORZ	r	Inclusive or to Register Zero	1	TPSU	v	Test Program Status, Upper, Selective	2
IORI,r	v	Inclusive or Immediate	2	TPSL	v	Test Program Status Lower, Selective	2
IORR,r	(*a)	Inclusive or Relative	2	INPUT/OUTPUT INSTRUCTIONS			
IORA,r	(*a),(X)	Inclusive or Absolute	3	WRD,r		Write Data	1
EORZ	r	Exclusive or to Register Zero	1	REDD,r		Read Data	1
EORI,r	v	Exclusive or Immediate	2	WRTC,r		Write Control	1
EORR,r	(*a)	Exclusive or Relative	2	REDC,r		Read Control	1
EORA,r	(*a),(X)	Exclusive or Absolute	3	WRTE,r	v	Write Extended	2
COMPARISON INSTRUCTIONS				REDE,r	v	Read Extended	2
COMZ	r	Compare to Register Zero	1	MISCELLANEOUS INSTRUCTIONS			
COMI,r	v	Compare Immediate	2	HALT		Halt, Enter Wait State	1
COMR,r	(*a)	Compare Relative	2	DAR,r		Decimal Adjust Register	1
COMA,r	(*a),(X)	Compare Absolute	3	TMI,r	v	Test Under Mask Immediate	2
ROTATE INSTRUCTIONS				NOP		No Operation	1
RRR,r		Rotate Register Right	1	<hr/>			
RRL,r		Rotate Register Left	1	¹ 2650B only			
BRANCH INSTRUCTIONS							
BCTR,v	(*a)	Branch on Condition True Relative	2				
BCFR,v	(*a)	Branch on Condition False Relative	2				
BCTA,v	(*a)	Branch on Condition True Absolute	3				
BCFA,v	(*a)	Branch on Condition False Absolute	3				
BRNR,r	(*a)	Branch on Register Non-Zero Relative	2				
BRNA,r	(*a)	Branch on Register Non-Zero Absolute	3				
BIRR,r	(*a)	Branch on Incrementing Register Relative	2				
BIRA,r	(*a)	Branch on Incrementing Register Absolute	3				
BDRR,r	(*a)	Branch on Decrementing Register Relative	2				
BDRA,r	(*a)	Branch on Decrementing Register Absolute	3				
BXA	(*a),(x)	Branch Indexed Absolute, Unconditional	3				
ZBRR	(*a)	Zero Branch Relative, Unconditional	2				

APPENDIX C

NOTES ABOUT THE 2650 MICROPROCESSOR

1. AUTO-INCREMENT, DECREMENT of index register. This feature is optional on any instruction which used indexing with the exception of BXA and BSXA. The increment or decrement occurs before the index register is added to the displacement in the instruction.
2. The contents of registers when used for indexing are considered to be unsigned absolute numbers. Consequently, index registers can contain values from 0 to 255. They "wrap-around" so that the number following 255 is 0.
3. Only absolute addressing instructions can be indexed.
4. The Branch on Incrementing Register or Decrementing Register instructions perform the increment or decrement before testing for zero. The only time the branch is not made is when the register contains zero.
5. All hardware relative addressing is implemented as modulo 8K and therefore relative addressing across the top of a page boundary will result in a physical address near the bottom of the page being accessed. For example:

1FFC₁₆ LODR,R2 \$+16

This instruction results, during execution, in accessing the byte at location 000C in the same

page as the instruction. Similarly, negative relative addresses from near the bottom of a page may result in an effective address near the top of the page.

6. Page boundaries cannot be indexed across.
7. Data can always be accessed across a page boundary through use of relative indirect or absolute indirect addressing modes.
8. The only way to transfer control to a program in some other page is to branch absolute or branch indirectly to the new page. Program execution cannot flow across a page boundary.
9. Unconditional branch or branch to subroutine instructions are coded by specifying a value of 3 in the register/value field of BSTA, BSTR, BCTA or BCTR. Example:

UN	EQU	3
	•••	
	•••	
	•••	
	BSTA,UN	PAL
	BCTR,3	LOOP

Unconditional branches or subroutine branches on condition false (BCFA, BCFR, BSFA, BSFR) are not allowed.

APPENDIX D

ASCII CODE

This table presents the only characters that the assembler will recognize in an A type constant and their equivalent codes in hexadecimal.

VALID CHARACTERS	ASCII CODE	VALID CHARACTERS	ASCII CODE
0	30	V	56
1	31	W	57
2	32	X	58
3	33	Y	59
4	34	Z	5A
5	35	blank	20
6	36	.	2E
7	37	(28
8	38	+	2B
9	39		7C
A	41	&	26
B	42	!	21
C	43	\$	24
D	44	*	2A
E	45)	29
F	46	;	3B
G	47	¬ or ~	7E*
H	48	-	2D
I	49	/	2F
J	4A	,	2C
K	4B	%	25
L	4C	- or ←	5F*
M	4D	>	3E
N	4E	?	3F
O	4F	:	3A
P	50	#	23
Q	51	@	40
R	52	.	27
S	53	=	3D
T	54	"	22
U	55	<	3C

*may have different graphic symbols on different output devices.



APPENDIX E

ASCII CHARACTER SET

ASCII CHARACTER SET (7-BIT CODE)									
L.S. CHAR	M.S. CHAR	0 000	1 001	2 010	3 011	4 100	5 101	6 110	7 111
0	0000	NUL	DLE	SP	0	@	P	`	p
1	0001	SOH	DC1	!	1	A	Q	a	q
2	0010	STX	DC2	"	2	B	R	b	r
3	0011	ETX	DC3	#	3	C	S	c	s
4	0100	EOT	DC4	\$	4	D	T	d	t
5	0101	ENQ	NAK	%	5	E	U	e	u
6	0110	ACK	SYN	&	6	F	V	f	v
7	0111	BEL	ETB	'	7	G	W	g	w
8	1000	BS	CAN	(8	H	X	h	x
9	1001	HT	EM)	9	I	Y	i	y
A	1010	LF	SUB	*	:	J	Z	j	z
B	1011	VT	ESC	+	;	K	[k	{
C	1100	FF	FS	,	<	L	\	l	
D	1101	CR	GS	-	=	M]	m	}
E	1110	SO	RS	•	>	N	↑	n	~
F	1111	SI	US	/	?	O	← or —	o	DEL

APPENDIX F POWERS OF TWO TABLE

2^n	n	2^{-n}
1	0	1.0
2	1	0.5
4	2	0.25
8	3	0.125
16	4	0.062 5
32	5	0.031 25
64	6	0.015 625
128	7	0.007 812 5
256	8	0.003 906 25
512	9	0.001 953 125
1 024	10	0.000 976 562 5
2 048	11	0.000 488 281 25
4 096	12	0.000 244 140 625
8 192	13	0.000 122 070 312 5
16 384	14	0.000 061 035 156 25
32 768	15	0.000 030 517 578 125
65 536	16	0.000 015 258 789 062 5
131 072	17	0.000 007 629 394 531 25
262 144	18	0.000 003 814 697 265 625
524 288	19	0.000 001 907 348 632 812 5
1 048 576	20	0.000 000 953 674 316 406 25
2 097 152	21	0.000 000 476 837 158 203 125
4 194 304	22	0.000 000 238 418 579 101 562 5
8 388 608	23	0.000 000 119 209 289 550 781 25
16 777 216	24	0.000 000 059 604 644 775 390 625
33 554 432	25	0.000 000 029 802 322 387 695 312 5
67 108 864	26	0.000 000 014 901 161 193 847 656 25
134 217 728	27	0.000 000 007 450 580 596 923 828 125
268 435 456	28	0.000 000 003 725 290 298 461 914 062 5
536 870 912	29	0.000 000 001 862 645 149 230 957 031 45
1 073 741 824	30	0.000 000 000 931 322 574 615 478 515 625
2 147 483 648	31	0.000 000 000 465 661 287 307 739 257 812 5
4 294 967 296	32	0.000 000 000 232 830 643 653 869 628 906 25
8 589 934 592	33	0.000 000 000 116 415 321 826 934 814 453 125
17 179 869 184	34	0.000 000 000 058 207 660 913 467 407 226 562 5
34 359 738 368	35	0.000 000 000 029 103 830 456 733 703 613 281 25
68 719 476 736	36	0.000 000 000 014 551 915 228 366 851 806 640 625
137 438 953 472	37	0.000 000 000 007 275 957 614 183 425 903 320 312 5
274 877 906 944	38	0.000 000 000 003 637 978 807 091 712 951 660 156 25
549 755 813 888	39	0.000 000 000 001 818 989 403 545 856 475 830 078 125
1 099 511 627 776	40	0.000 000 000 000 909 494 701 772 928 237 915 039 062 5

APPENDIX G

HEXADECIMAL-DECIMAL CONVERSION TABLES

From hex: locate each hex digit in its corresponding column position and note the decimal equivalents. Add these to obtain the decimal value.

From decimal: (1) locate the largest decimal value in the table that will fit into the decimal number to be converted, and (2) note its hex equivalent and hex column position. (3) Find the decimal remainder. Repeat the process on this and subsequent remainders.

The table on pages 85-88 provides for direct conversion of hexadecimal and decimal numbers in these ranges:

Hexadecimal	Decimal
000 to FFF	0000 to 4095

In the table, the decimal value appears at the intersection of the row representing the most significant hexadecimal digits (16^2 and 16^1) and the column representing the least significant hexadecimal digit (16^0).

Example: $C21_{16} = 3105_{10}$

HEX	0	1	2
C0	3072	3073	3074
C1	3088	3089	3090
C2	3104	3105	3106
C3	3120	3121	3122

HEXADECIMAL COLUMNS											
6		5		4		3		2		1	
HEX = DEC		HEX = DEC		HEX = DEC		HEX = DEC		HEX = DEC		HEX = DEC	
0	0	0	0	0	0	0	0	0	0	0	0
1	1,048,576	1	65,536	1	4,096	1	256	1	16	1	1
2	2,097,152	2	131,072	2	8,192	2	512	2	32	2	2
3	3,145,728	3	196,608	3	12,288	3	768	3	48	3	3
4	4,194,304	4	262,144	4	16,384	4	1,024	4	64	4	4
5	5,242,880	5	327,680	5	20,480	5	1,280	5	80	5	5
6	6,291,456	6	393,216	6	24,576	6	1,536	6	96	6	6
7	7,340,032	7	458,752	7	28,672	7	1,792	7	112	7	7
8	8,388,608	8	524,288	8	32,768	8	2,048	8	128	8	8
9	9,437,184	9	589,824	9	36,864	9	2,304	9	144	9	9
A	10,485,760	A	655,360	A	40,960	A	2,560	A	160	A	10
B	11,534,336	B	720,896	B	45,056	B	2,816	B	176	B	11
C	12,582,912	C	786,432	C	49,152	C	3,072	C	192	C	12
D	13,631,488	D	851,968	D	53,248	D	3,328	D	208	D	13
E	14,680,064	E	917,504	E	57,344	E	3,584	E	224	E	14
F	15,728,640	F	983,040	F	61,440	F	3,840	F	240	F	15
BYTE 1				BYTE 2				BYTE 3			

COMMAND/DIRECTIVE SUMMARY

Command/Directive	Description	Page Number
{name} ACON expression	Assembler allocates two successive bytes	27
.ASSIGN &<variable> = <expression>	Value of expression assigned to variable	59
CEJE expression	Same as EJE if remaining number of lines is less than value of expression	33
label CSECT	Identify an object module	43
{name} DATA argument	Assembler allocates exact number of bytes	28
.DECLARE &<variable>,...	Declares variables for use by macroprocessor	59
EJE	Assembler advances listing to top of page	32
ELSE	Generate different sequences of statements	40
END {expression}	Tells assembler last statement has been input	31
ENDIF	End of conditional assembly construct	40
ENTRY label list ...	Specify symbols referenced by other programs	44
name EQU expression	Equate symbol in label field with evaluable expression in operand field	25
EXTRN label list ...	Specify symbols defined in some other program which are referred to in this program	45
.GENERATE &<variable> = <expression1> <expression2> .END	Statement pair delimits group of statements to be processed iteratively	59
{name} IF expression	Programmer can vary the sequence of statements	40
.IF <expression> .&<macrolabel>	Conditionally include block of text	59
_INCLUDE name	Include separately assembled object program module	49
.JUMP .&<macrolabel>	Indicate next statement to be scanned	59
LINK <input>, [<loadmap>], [<bfile>]	Link editor command	46
MAC <source>, <expanded source>, [<macro file>], [<error list>]	Macro preprocessor command	50
.MACRO ⋮ .MEND } macro definition	Command pair brackets each macro definition	60
.&<macrolabel>	Label macro command	59
.MESSAGE <text>	Display text on system console	60
.NOP	No op used for defining macrolabels	60
{name} ORG expression	Set assembly time location counter	24
PCH { ON } { OFF }	Resume/discontinue output of object program module	37
_PHASE <bfile> [<origin>] [<ID>]	Name a phase to be created by link editor	48
PRT { ON } { OFF } { GEN } { NOGEN }	Resume/discontinue printing of assembled program	34
RASM <ifile>, [<list>], [<ofile>], [<print>]	Relocatable assembler command	18
REPRO	Copy link editor commands from source module	38
{name} RES expression	Assembler reserves contiguous bytes of storage	30
name SET expression	Equate symbol in label field with expression in argument field	26
SPC expression	Assembler skips or spaces a number of lines	35
.SUBSTRING &<variable> = &<variable> (<expression1>, <expression2>)	Second variable substring assigned to first variable	60
.TAB character	Sets the tab character	60
TITL expression	Specifies a title for the assembly listing	36

COMMENT SHEET

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This form is not intended to be used as an order blank. Signetics Corporation solicits your comments about this manual with a view to improving its usefulness in later editions.

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